

# The Design of Network Service Architecture and Its Application to AG 2



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# Outline

- Review:
  - Assumptions of AG
  - Goals of AG2
- Novel proposal for Network Service Architecture (NSA)
  - Properties of NSA
  - General approach of NSA
- Current and future work
- Summary



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# Assumptions of AG

- The original assumption:
  - Higher bandwidth;
  - Higher computational capability;
- Implementation of AG1
  - Location constraints;
  - Resource constraints (bandwidth, computing power, etc.).



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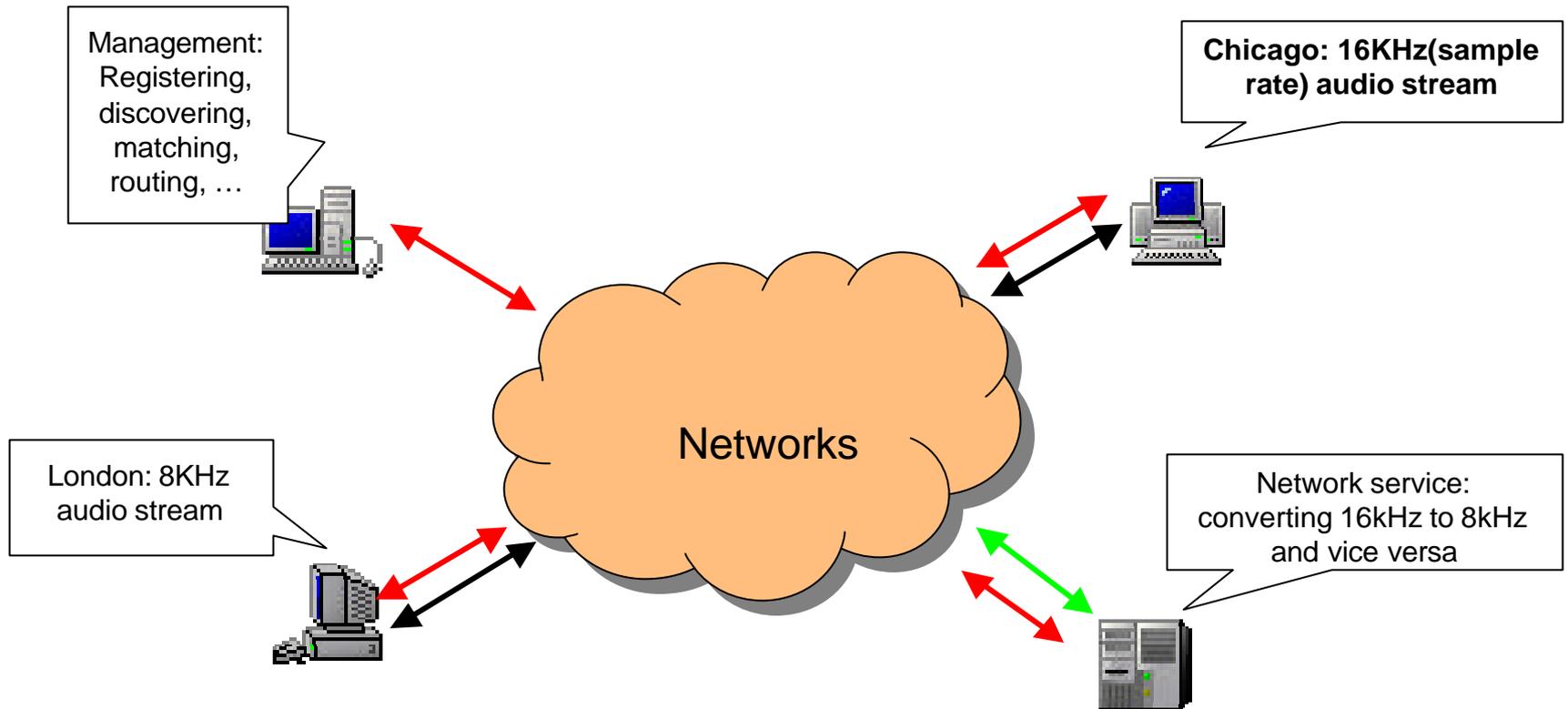
# Goals of AG2

- Building environment for better remote interactions
  - Free of location constraints
- Relaxing resource constraints
  - Various bandwidth, computing capability
  - Multi-platform support



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# Example

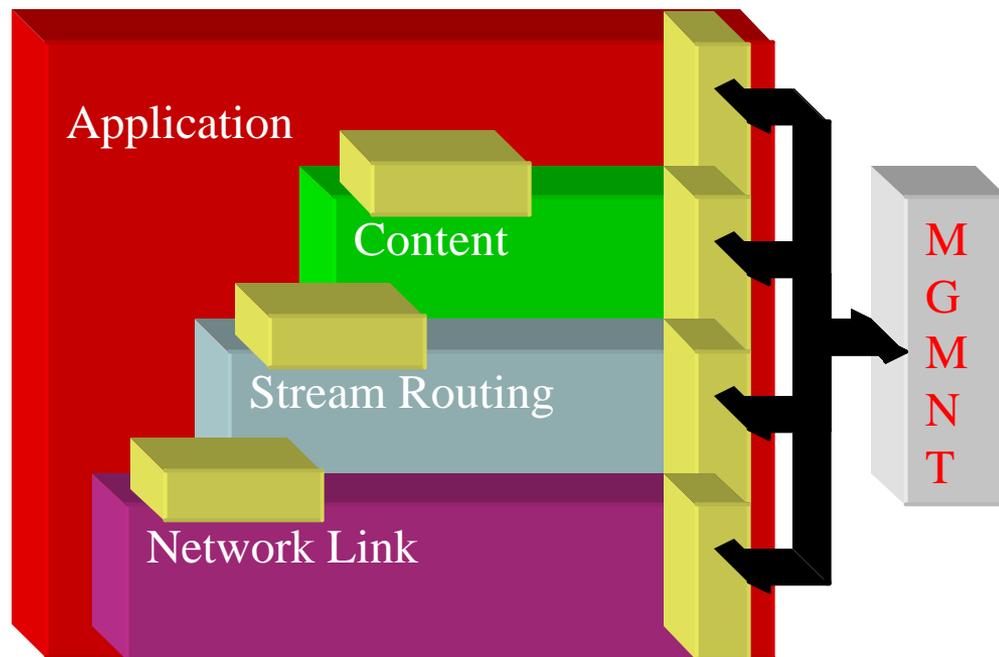


# Properties of Network Service and Network Service Architecture

- Properties of network service
  - Network service as a network resource object
  - Decentralization
  - Heterogeneous conditions
- Properties of network service architecture
  - Soft-state
  - Decentralization
  - Self-adaptable
  - Transparent
  - Stream-Capability



# Network Service Architecture

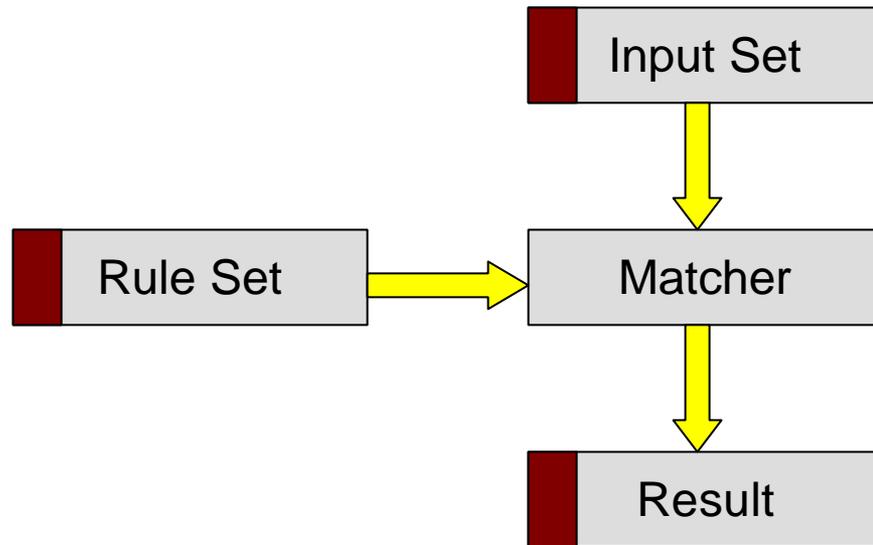


# Network Service Architecture

- Applications
  - Resource, networks, ...
- Content-Interceptor
  - description files defined by specific schemas
- Stream Route
  - Finding optimal streaming topology
- Network Links
  - Reflector, Bridges, ...



# General Approach: General Matcher



# General Approach: Matching Algorithm

- Capability Space and Capability Vector:

$$\mathbf{p} = (v_1, v_2, \dots, v_n) \in R^n$$

- Network Service as Transformation Matrix:

$$\mathbf{v}_1 = T\mathbf{v}_2$$

# General Approach: Matching Algorithm (cont.)

- Finding the best service in two sets:

$$T_{1,2} = T(S_1^n, S_2^n) = [T_{i,j}]_{\max_{\|\bar{v}_{1,i}\|}}, \quad \text{where } e_{\bar{v}_{1,i}}^\omega = T_{i,j}^\omega \bar{v}_{2,j}$$

- Finding a set of services for multi-sets:

$$S_1^n, S_2^n, \dots, S_k^n, S_i^n \subseteq R^n, i = 1, \dots, k.$$

$$S_1^{\cup n} \perp S_2^{\cup n} \perp \dots \perp S_l^{\cup n}, l \leq k$$

$$\{T_{1,2}, T_{1,3}, \dots, T_{i,j}, \dots, T_{l-1,l}\}, \text{ where } T_{i,j} = T(S_i^{\cup n}, S_j^{\cup n})$$

# General Approach: Stream Topology Algorithm

- Active unconnected session:
- We borrow nodes from non-active session to comprise a connected graph:
- Finding the minimum cost connected graph:

$$G(V, E)$$

$$G''(V'', E'')$$

$$C_{opt} = \min_{V'', E''} (C(G''(V'', E''))) )$$

# Current and Future Work

- Current work
  - Simulation of a model management module: two inputs → one best resolution choice or a list of choices
- Future work
  - Implementation of general matcher
  - Implementation of stream topology algorithm
  - Integration of network service architecture



# Summary

- Novel proposal of Network Service Architecture
  - General Matcher and Matching Algorithms
  - Streaming Topology Algorithms

